

# A Robust Framework for Socially Responsive Services: A Constraint-Based Social Network Representation



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## 1 Introduction

NGOs can be defined as non-profit organizations and a group of members who work in some ways for serving the individuals or a community [1]. A well-established literature shows that the service provider and the service receiver influence the behavior of each other and which is highly influenced by social network structure. Therefore, NGOs can be defined as private organizations who work in local level or national level or international levels whose objective is to help or share the resources among needy parts of the society. It is a non-profit community-based organization.

The main goal of NGOs is to share resources [2]. Here resource signifies any types of help that can be provided by the individuals or community. It may be idea, some solution, physical resource, or some kinds of interdependency, and it will satisfy the demands of the user. It should be feasible in true sense; that is, it should

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be technologically, economically, and culturally feasible [3]. Here our main focus is that the distribution should be fair and unbiased. The needy people should get the help. The help may be educational like a fellowship for a poor student, medical help for a patient, or some help for a disabled [4]. Otherwise, we may face the following major problems:

- Overall crisis.
- Requirement of service of needy people.
- Development of poorer section of society.

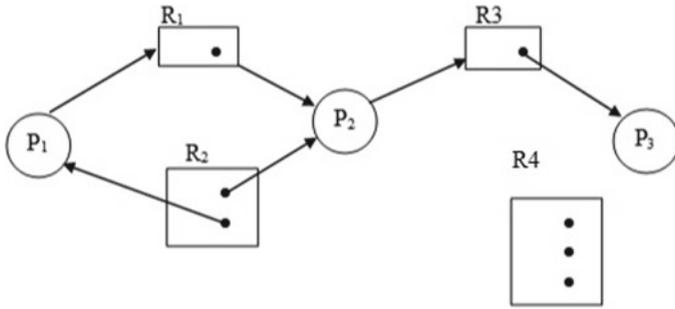
Galvin et al. [5] propose the concept of resource management and the usage of resources. Resources will be supplied by the assigned agency. Better collaboration and interaction are possible with the advancement of Internet with time. The social network paved its way into popularity by superseding the Internet. Social network analysis is a vast area of research which deals with the understanding and analysis of the behavior of individuals or community. It predicts the prospective existence of the relationships between community members.

Socially Responsive Resource Usage Protocol (SRRUP) defined to standardized social ethics and utilization of resources [6, 7]. This protocol helps to emerge in an improvised environment by providing the way of accessing the resources to needy people. To allocate the resources, more socially distant individual is given more priority than less socially distant individual. Social distance of the user represents the defined logical distance between user and resource [8–10]. A resource has three parameters: place, cost, and constraints, causing prevailing consequences in its usages. Place can be defined where the resource can be availed [11–13]. A person who needs resources pay the price to avail the resources is known as cost. A set of affirmation and considerations measure to a resource usage are the constraints.

In our proposed framework, we have made the following contributions:

- The distribution of resources will be impartial and those resources can be avail for those people who genuinely have the requirements.
- The resource will be available to the persons based on their financial, social, disadvantageous positions, and their locations (remote or rural).

In operating system, we termed user as a process. Process enters in waiting state if it does not get requested resources. Some time waiting, process does not come out from waiting state because requested resources are held by some other waiting process and so on then it becomes deadlock. Such situation occurs for inquiry table distribution of resources. For this work, we have given attention to fair resource allocation. To describe it precisely, best option is to use a directed graph; this is called resource distribution graph. The graph has nodes  $V$  and edges  $E$ .  $V$  is partitioned into  $\{P_1, P_2, \dots, P_n\}$ , the set of active processes and  $R = \{R_1, R_2, \dots, R_m\}$ , the set resource types in the system. Figure 1 depicts the graph where a directed edge from process  $P_i$  to resource type  $R_j$  is referred by  $P_i \rightarrow R_j$ ; it shows that process  $P_i$  has requested an instance resource type  $R_j$  and is currently waiting for that resource. A directed edge from resource type  $R_j$  to process  $P_i$  is referred by  $R_j \rightarrow P_i$ ; it means that an instance of resource type  $R_j$  has been allocated to process



**Fig. 1** Resource allocation graph

$P_i$ . A directed edge  $P_i \rightarrow R_j$  is named a request edge; a directed edge  $R_j \rightarrow P_i$  is named an assignment edge.

Though different resource allocation graph is used in many applications like job processing, hospital management, and real-time scheduling, this is mainly used for avoidance of deadlock where same resources are demanded by multiple processes and others are not releasing the resources [14]. In our proposed research work, multiple instances of resource are possible. Hence, this situation will not raise.

Section 2 presents the working mechanism of the proposed model. The experimental result and comparison studies are presented in Sect. 3. In the last section, future work and conclusions has been provided.

## 2 Proposed Model

### 2.1 Concept of the Proposed Model

In our earlier work [1], if a volunteer is not able to provide a support due to lack of resources then he will wait for required resources. To improve the performance by reducing the waiting time of our earlier work, we have proposed a new scheme.

If two volunteers have the less resources than receiver’s requirement like volunteers  $v_1 < r$  and  $v_2 < r$ . If  $v_1 + v_2 \geq r$  then both volunteers jointly provide resources to that receiver for completion of his job. This way it is reducing the waiting time and utilizing the resources effectively.

In this model, we have designed a network comprised of two layers: one layer represents the service provided by the service provider and other layer represents the members who will receive the services. The services will be received by the

weaker sections of the society such as orphans, disabled, poor or needy, and old peoples. The service will be given voluntarily by individual or by a community. Here the volunteers mainly provide two types of services. One is to identify the deserving persons who are willing to take services. Second responsibility is to accumulate the resources to provide services without interruption to the needy, old, disabled, orphan, or persons staying in remote areas. The services required vary from user to user. Some users may require financial help for continuing their study. Some may require physical assistance or medical supervision for their agony and pain. Volunteer also has different resources with different capacities. At a time, more than one service can be provided by volunteer. If a user requires services a volunteer alone is not able to help, then services can be provided by more than one volunteer to the same user.

So the communication is one-to-one (volunteer to receiver), one-to-many (receiver to volunteers), and many-to-one (receivers to volunteer). The volunteers may approach external sources to serve the users because requirements of the user vary. For approaching the external sources, volunteers have to follow system rules: Volunteer cannot take more than user's requirement. Based on resource availability the needed service will be given. Volunteers of a specific group of users communicate with other volunteer if required. User quits the organization or concern volunteer may stop his support to user after fulfilling the requirement.

## ***2.2 Mathematical Representation of the Proposed Model***

We have introduced user by  $u$  and volunteers by  $v$  with appropriate suffixes. Users seek their requirements through  $u(x_1, x_2, \dots, x_k)$  and volunteer represents services through  $v(y_1, y_2, \dots, y_p)$  where  $x_i$  and  $y_i$  represent direct amount that should be non-negative real numbers. It is used to find the strength of the organization on the basis of user need and volunteer service.

Using community detection in social network analysis [5, 6], we organize and find various categories of person with different need and support to run the organization smoothly.

We prefer two types of sorting:

Case I. User-based category.

Case II. Volunteer-based category.

This can be classified to get a real scenario.

**Case I** Suppose that more than one service can be provided by each volunteer but he has only one service to provide. Then we seize an estimable amount of users to him.

Let  $v_1(y_1)$  represent a volunteer with a type of support  $s_1$  (like disables, orphans).

Let a category of users  $\{u_{11}, u_{12}, \dots, u_{1k}\}$  are associated with  $v_1$  with demand of  $y_{11}, y_{12}, \dots, y_{1k}$  where  $u_{1i}(y_{1i})$  represents a typical member of this category.

Now there are two possibilities:

**Lemma 1**  $s_1 \geq \sum_{j=1}^k y_{1j}$

(Sum of required is less than service quantum)

Here, either the volunteer or organization  $v_1$  themselves asks for more users to extend this task.

Now we can extend  $u_{1k+1}(y_{1k+1}), u_{1k+2}(y_{1k+2}), \dots, u_{1p}(y_{1p})$  where  $p > k$  and  $s_1 \geq \sum_{j=1}^p y_{1j}$ .

This process goes on till  $\sum_{j=1}^p y_{1j} \leq s_1 < \sum_{j=1}^{p+1} y_{1j}$ .

This improve the performance of  $v_1$ . Improvement of  $v_1$  is evaluated on the basis of more users ( $u_{1j}$ 's).

**Lemma 2**  $s_1 \geq \sum_{j=1}^k y_{1j}$

In this case, the support of  $v_1$  is not sufficient to satisfy the needs of  $u_{1k}$ 's. So either the volunteer himself ( $v_1$ ) or NGO contacts some other volunteer  $v_2$  of similar task as demanded by  $u_{1k}$ 's. Now here options are:

(a) if  $\sum_{j=1}^l y_{1j} \leq s_1 < \sum_{j=1}^k y_{1j}$  where  $l < k$ , we can detach those  $\{u_{1l+1}, u_{1l+2}, \dots, u_{1k}\}$  from  $v_1$  and associate them to  $v_2$  directly. Here  $v_2$  is an existing volunteer of NGO.

or (b)  $v_2$  cannot approach directly to  $u_1$ 's. NGO will link up  $v_2$  to  $v_1$  so now:  $s_1 + s_2 \geq \sum_{j=1}^k y_{1j}$

where  $s_1$  and  $s_2$  are the support of  $v_1$  and  $v_2$ , respectively. In this case, for purely temporary basis either NGO or volunteer gets a support from outside.

**Case II** If each  $u_i$  needs many services and no one volunteer alone is sufficient to fulfill their needs, then each  $u$  will be assigned to a group of  $v$ 's. Here user

represents his needs through  $u_1(y_{11}, y_{12}, \dots, y_{1k})$  and group of volunteer represent their strength through  $\{v_{11}(s_{11}), v_{12}(s_{12}), \dots, v_{1k}(s_{1k})\}$ . We expect  $y_{1j} \leq s_{1j}$ , where  $y = 1, 2, \dots, k$  else NGO has to include more  $v$ 's to fulfill the needs of  $u_1$ .

**Lemma 1** Suppose  $y_{1j} = s_{1j} \forall j = 1, 2, \dots, k$

Then all the needs of  $u_1$  are fulfilled and  $u_1$  may leave.

**Lemma 2** Suppose  $y_{1j} < s_{1j}$  for some  $y = p \leq k$ .

Then the  $v_{1p}(s_{1p})$  can support some other users 'u' also. NGO will utilize  $v_{1p}(s_{1p})$  remaining services to fulfill the needs of other user  $u_2$  like  $u_2(y_{2p})$  for which  $y_{1p} + y_{2p} \leq s_{1p}$ .

This way such  $v_{1p}$  are associated with more users. This process repeats until  $\sum_{j=1}^k y_{1j} \leq s_{1j}$ .

Here, the public presentation of the organization possibly is assessed by number of users associated with a volunteer and the number of users leaving the system with satisfaction.

### 2.3 Algorithms of the Proposed Model

Here, we declare all attributes which exist in this proposed algorithm.

$u \leftarrow$  User.

$y \leftarrow$  Requirement.

$u_{1j} \leftarrow y_{1j}$  Users ( $u_{11}, u_{12}, \dots, u_{1j}$ ) with requirements ( $y_{11}, y_{12}, \dots, y_{1j}$ ).

$v \leftarrow$  Volunteer.

$s \leftarrow$  Service.

$v \leftarrow s$  Volunteer ( $v$ ) provides service ( $s$ ).

$v_1 \leftarrow s_1, v_1$  provides  $s_1$ -type service,  $v_2 \leftarrow s_2, v_2$  provides  $s_2$ -type service etc.

In this algorithm, our aim is to provide services ( $s$ ) to a user ( $u$ ), who are registered ( $u \leftarrow y$ ) with requirements ( $y$ ). Volunteer ( $v$ ) provides service to user ( $u$ ) to fulfill their needs. In the first case, the organization is robust, while in the second case the organization asks for improvement.

Algorithm 1: Algorithm for Case I

```

Begin
 $u_{1j} \leftarrow r_{1j}, \forall j=1, 2, \dots, k;$ 
 $v_1 \leftarrow s_1, v_2 \leftarrow s_2$ 
if  $s_1 > r_{11}$ ,
    | then  $v_1$  supports  $r_{11}$ 
end
for  $l=2$  to  $k$  do
    | if  $s_1 > \sum_{j=1}^l r_{1j}$  then
    | |  $v_1$  supports  $\{x_{11}, x_{12}, \dots, x_{1l}\}$ 
    | | end
    | if  $\sum_{j=1}^l r_{1j} \leq s_1 \leq \sum_{j=1}^{l+1} r_{1j}$  then
    | |  $v_1$  supports  $\{x_{11}, x_{12}, \dots, x_{1l}\}$  but not  $x_{1l+1}$ 
    | | end
end
select  $v_2(s_2)$  such that
 $s_1 + s_2 \geq \sum_{j=1}^k r_{1j}$  or  $s_2 \geq \sum_{j=1}^k r_{1j} - s_1$  then
 $\{v_1, v_2\}$  supports  $\{x_{11}, x_{12}, \dots, x_{1k}\}$ 
End
    
```

Algorithm 2: Algorithm for Case II

```

Begin
 $u_1 \leftarrow \{r_{11}, r_{12}, \dots, r_{1k}\};$ 
for  $l=1$  to  $k$  do
    |  $v_1 \leftarrow s_{1l}$ 
end
for  $l=1$  to  $k$  do
    | if  $r_{1l} \leq s_{1l}$  then
    | |  $u_1$  leaves the system as satisfied customer
    | |  $\{v_{11}, v_{12}, \dots, v_{1k}\}$  supports  $u_1$ .
    | | end
end
Now  $u_2 \leftarrow \{r_{21}, r_{22}, \dots, r_{2k}\}$ 
for  $l=1$  to  $k$  do
    | if  $r_{2l} \leq s_{1l} - r_{1l}$ 
    | | then  $u_2$  leaves the system as satisfied customer
    | |  $\{v_{11}, v_{12}, \dots, v_{1k}\}$  supports  $\{u_1, u_2\}$ 
    | | end
end
for  $p=1$  to  $k$  do
    | if  $r_{2p} \leq s_{1p} - r_{1p}$ 
    | | then some  $v'_{1p}$  enter the system with  $v'_{1p} \leftarrow s'_{1p}$  and
    | |  $r_{2p} \leq s'_{1p} + s_{1p} - r_{1p}$ 
    | |  $u_2$  leaves the system as satisfied manner.
    | |  $\{v_{11}, v_{12}, \dots, v_{1k}, v'_{1p}\}$  supports  $\{u_1, u_2\}$ 
    | | end
end
    
```

This process continues until all  $u_i$ 's and  $v_k$ 's are exhausted.

End

### 3 Experimental Result Discussion

Here, we elaborated our program output. We have taken all possible kinds of inputs and elaborated desired outputs of concerned inputs.

#### 3.1 Experimental Environments

We have implemented these algorithms and tested using all possible desired inputs using a Java program for the preliminary study and executed it on a Lenovo Miix 510-121SK laptop with 2.3 GHz Intel Core i5-6200U 6th Gen processor, 8 GB DDR3 RAM and 256 GB hard disk, Windows 10 with 64-bit operating system and Java1.8. In this program, our objective is to make a communication among various methods. Program works on the basis of client and server program. Client represents the user and server represents the providers. There are three types of user-defined communication methods:

- Volunteer\_to\_receiver—This method represents the task like one volunteer is alone quite sufficient to fulfill the needs of a user of all type of services.
- Receiver\_to\_volunteers—This method provides the dependency of volunteers to fulfill the requirements of receiver.
- Receivers\_to\_volunteers—This method works as services of more than one user can be fulfilled by more than one volunteer. Above communications can be initialized in various ways like users represented by ‘ $u$ ’ and volunteers by ‘ $v$ ’. The needs of user are represented by  $y_i, i = 1, 2, \dots, n$  and services of volunteers are represented by  $s_j, j = 1, 2, \dots, n$ . Table 1 represents the input provided to the system and Table 2 reflects the desired output.

**Table 1** User needs table

User/need	$r_1$	$r_2$	$r_3$	...	...	$r_n$
$u^1_{1000}$	1	0	0	0	0	0
$u^1_{0100}$	0	1	0	0	0	0
...	...	...	...	...	...	...
$u^2_{1100}$	1	1	0	0	0	0
$u^2_{1010}$	1	0	1	0	0	0
...	...	...	...	...	...	...
$u^3_{1110}$	1	1	1	0	0	0
...	...	...	...	...	...	...
Sum	=4	=3	=	=	=	=
Total =						

**Table 2** Volunteer support table

Provider/support	$s_1$	$s_2$	$s_3$	...	...	$s_n$
$v_1^1$	$k_1$	0	0	0	0	0
$v_2^1$	0	$k_2$	0	0	0	0
...	...	...	...	...	...	...
$v_1^2 2^2$	$p_1$	$p_2$	0	0	0	0
$v_1^3 3^2$	$p_1$	0	$p_3$	0	0	0
...	...	...	...	...	...	...
$v_1^5 2^4 3^3$	$q_1$	$q_2$	$q_3$	0	0	0
...	...	...	...	...	...	...
Sum	=	=	=	=	=	=

**Notation:**  $u^i$  is a user ( $i = 1$  to  $n$ )

$u^1_{1000}$  denotes user wants single resource of  $y_1$ .

$u^2_{0100}$  denotes user wants single resource of  $y_2$ .

$u^1_{1100} = u^2_{1100}$  denotes user wants two resources  $y_1$  and  $y_2$ , respectively.

$u^1_{1110} = u^2_{1110} = u^3_{1110} = \dots$  represents a user needs three resources  $y_1, y_2,$  and  $y_3$ .

$u^1_{1111} = u^2_{1111} = u^3_{1111} = u^4_{1111} \dots$  represents a user with four resources  $y_1, y_2, y_3,$  and  $y_4$ .

$v_1^1$  has a strength of single resource  $r_1$  for type  $s_1$ .

$v_2^1$  has a strength of single resource  $r_2$  for type  $s_2$ .

$v_1^{526}$  has a strength of five resources  $r_1$  for type  $s_1$  and six resources  $r_3$  for type  $s_3$ .

.....

$v_i^j$  has a strength of  $j$  for type  $s_i$  where  $i$  represents types of service and  $j$  represents for some resources.

**Simple Case:** Assume each user ‘ $u$ ’ needs one resource of size ‘1’ only. That is,  $y_1 = y_2 = \dots = y_n = 1$ , users with variable requirements, e.g.,  $u^1_{1000}(y_1) y_1 \geq 1$ .

### 3.2 Result of Various Communication Method

- In this function, we have supplied different kinds of inputs for all the above-mentioned methods. For volunteer\_to\_receiver communication method, we represent through Table 3. Here  $u^1_{3000}, u^2_{0500},$  and  $u^3_{0020}$  denotes corresponding user and  $v_1^5, v_1^{526},$  and  $v_3^8$  the volunteer. Here  $u^1$  wants three resources of type  $s_1, u^2$  wants five resources of type  $s_2,$  and  $u^3$  requires two resources of type  $s_3,$  while (Table 4)  $v_1, v_2,$  and  $v_3$  contain 5, (5 of  $s_1, 6$  of  $s_2$ ),

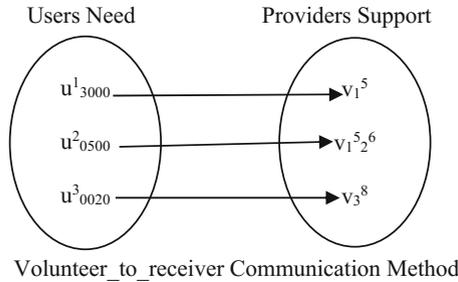
**Table 3** User needs table of receiver

User/need	$r_1$	$r_2$	$r_3$	$r_4$
$u^1_{3000}$	3	0	0	0
$u^2_{0500}$	0	5	0	0
$u^3_{0020}$	0	0	2	0

**Table 4** Provider support table of volunteer

Provider/support	$s_1$	$s_2$	$s_3$
$v_1^5$	5	0	0
$v_1^5, 2^6$	5	6	0
$v_3^8$	0	0	8

and 8 resources, respectively. So,  $v_1$  supports three resources to  $u^1$ ,  $v_2$  supports five resources of  $s_2$  to  $u^2$ , and  $v_3$  gives two resources to  $u^3$ . Now,  $v_1$  has two,  $v_2$  has (5 of  $s_1$ , 1 of  $s_2$ ), and  $v_3$  has six resources.



- For receiver\_to\_volunteers, we represent through Table 5. Here  $u^4_{4020}$  is denoted as a user which needs four resources of service  $s_1$  and two resources of  $s_3$ . Since (Table 6)  $v_1$  alone is not able to fulfill user need, because it has only two resources of  $s_1$ , so  $v_1$  needs help from another volunteer to satisfy user. Here  $v_1$ , and  $v_2$  support two and two resources of type  $s_1$ , respectively, and  $v_3$  supports two resources of  $s_3$ . Now,  $v_1$  has exhausted,  $v_2$  has (3 of  $s_1$ , 1 of  $s_2$ ), and  $v_3$  has 4 of  $s_3$ -type resources.

**Table 5** User needs table of receiver\_to\_volunteers

User/requirement	$r_1$	$r_2$	$r_3$	$r_4$
$u^4_{4020}$	4	0	2	0

**Table 6** Provider support table of receiver\_to\_volunteers

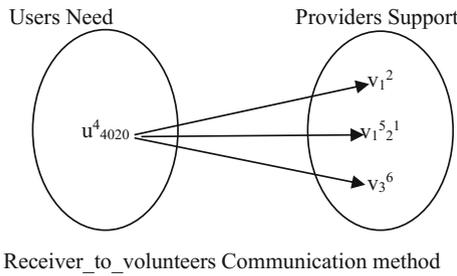
Provider/strength	$s_1$	$s_2$	$s_3$
$v_1^2$	2	0	0
$v_1^5, 2^1$	5	1	0
$v_3^6$	0	0	6

**Table 7** User needs table of receiver\_to\_volunteers

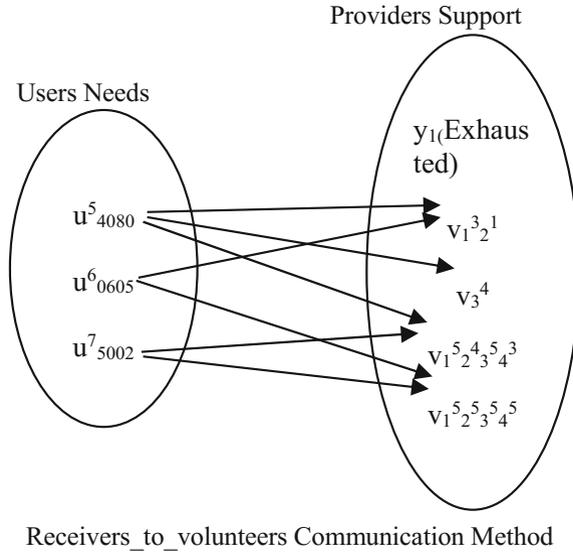
User/requirement	$r_1$	$r_2$	$r_3$	$r_4$
$u^5_{4080}$	4	0	8	0
$u^6_{0605}$	0	6	0	5
$u^7_{5002}$	5	0	0	2

**Table 8** Provider supports table of receivers\_to\_volunteers

Provider/strength	$s_1$	$s_2$	$s_3$	$s_4$
$v_1$	Exhausted	0	0	0
$v_1^3_2^1$	3	1	0	
$v_3^4$	0	0	4	0
$v_1^5_2^4_3^5_4^3$	5	4	5	3
$v_1^5_2^5_3^5_4^5$	5	5	5	5



- For receivers\_to\_volunteers, we represent through Table 7. Here,  $u^5_{4080}$ ,  $u^6_{0605}$ , and  $u^7_{5002}$  represent respective users. In Table 8, two more donors ( $v_1^5_2^4_3^5_4^3$  and  $v_1^5_2^5_3^5_4^5$ ) have been added from external sources. Here  $u^5_{4080}$  wants four resources of  $s_1$  and eight resources of  $s_3$ . From Table 8,  $v_1^3_2^1$  supports three resources of  $s_1$ ,  $v_1^5_2^4_3^5_4^3$  supports one resource of  $s_1$ , and  $v_3^4$  supports four resources of  $s_3$  and  $v_1^5_2^4_3^5_4^3$  supports four resources of  $s_3$ . Now,  $v_1^3_2^1$  has  $v_2^1$  only,  $v_3^6$  has exhausted, and  $v_1^5_2^4_3^5_4^3$  has  $v_1^5_2^4_3^5_4^3$ . User  $u^6_{0605}$  wants six resources of  $s_2$  and five resources of  $s_4$ , but  $v_2^1$  does not have sufficient resources to support. So it requires support from  $v_1^4_2^4_3^3_4^3$  and  $v_1^5_2^5_3^5_4^5$ . After giving the resources to  $u^6_{0605}$ ,  $v_2^1$  exhausted,  $v_1^4_2^4_3^3_4^3$  has  $v_1^4_3^1$  and  $v_1^5_2^5_3^5_4^5$  has  $v_1^5_2^4_3^5_4^3$ . User  $u^7_{5002}$  needs five resources of  $s_1$  and two resources of  $s_4$ .  $v_1^4_3^1$  supports four resources of  $s_1$  and  $v_1^5_2^4_3^5_4^3$  supports one resource of  $s_1$  and two resources of  $s_4$ . After giving the resources to  $u^7_{5002}$ ,  $v_1^4_3^1$  has  $v_3^1$  and  $v_1^5_2^4_3^5_4^3$  has  $v_1^4_2^4_3^5_4^1$ .



The number of satisfied user ( $u_i$ 's) and number of registering new users ( $u_i$ 's) measure the strength of the system.

### 4 Conclusions and Future Work

This framework suggests to provide services to various sections of the society and helps to improve the efficiency, effectiveness, and reliability of the system. It suggests couple of systematic methodologies like one volunteer to one receiver, one receiver to many volunteers, and many receivers to many volunteers to allocate the resources. In this chapter, we have formalized the holistic approach for fair distribution of service from one network to another network. Here one network represents the list of newly registered users and other represents the volunteers with types of services. Services have a broad range of latitude which leads the challenges for the system approaches to a service provider. There are two constraint matters to support a kind of services: identify the needy and supply the services to fulfill their requirements. For identification, either volunteer identifies or user has to approach to NGO. System provides the kind of services with the help of three communication methods. Once the user's needs are fulfilled, then either user might leave or volunteer stops his service. In both the cases, user has to give their feedback to the system. If the user leaves the system with satisfaction, then system is robust else system requires support. In implementation, we tried to assign providers to users on the basis of user's choice (preferential allocation in a prescribed order, say  $u_1$  to  $v_1$  like that).

In future work, user's feedback can be considered for better improvement of the system. We have to make sure best utilization of resources by keeping providers non-idle and consistent use of resources. We need to check the nearness between the resource providers. We should make the user as a temporary provider like increasing the demand of the user like  $u^1_{6000}$  which represents a user with need of  $y_1 = 6$ . Here  $u_1$  should be a supplier on some conditions but there is no provision for  $y_1$  in the existing system.

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